## CRDT Sets: Paper to Product

(Or Everything You Always Wanted to Know About ORSets\* (\*But Were Afraid to Ask))

## What?

- Why Riak?
- What is Riak?
- What's a CRDT, anyway?
- A replicated set

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# Why Riak?

## Scale Up

\$\$\$Big Iron (still fails)



### Scale Out

Commodity Servers CDNs, App servers Expertise



# Fault Tolerance

# Low Latency

### Low Latency

Amazon found every 100ms of latency cost them 1% in sales.



## Low Latency

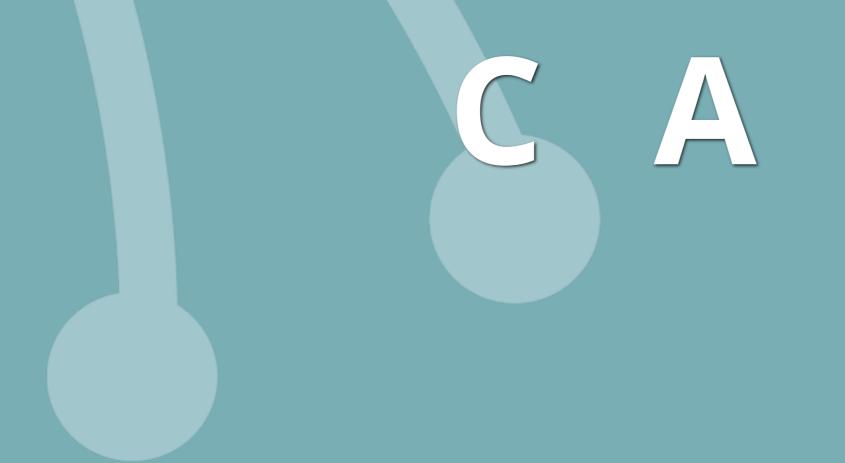
Google found an extra 0.5 seconds in search page generation time dropped traffic by 20%.

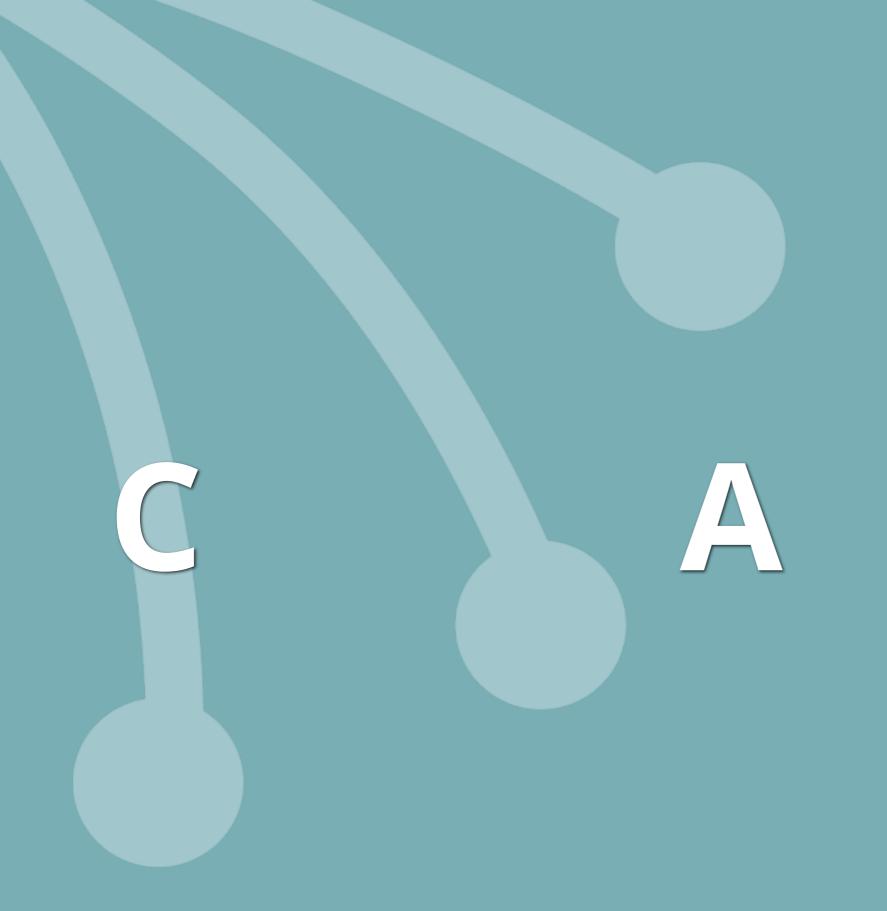


Trade Off

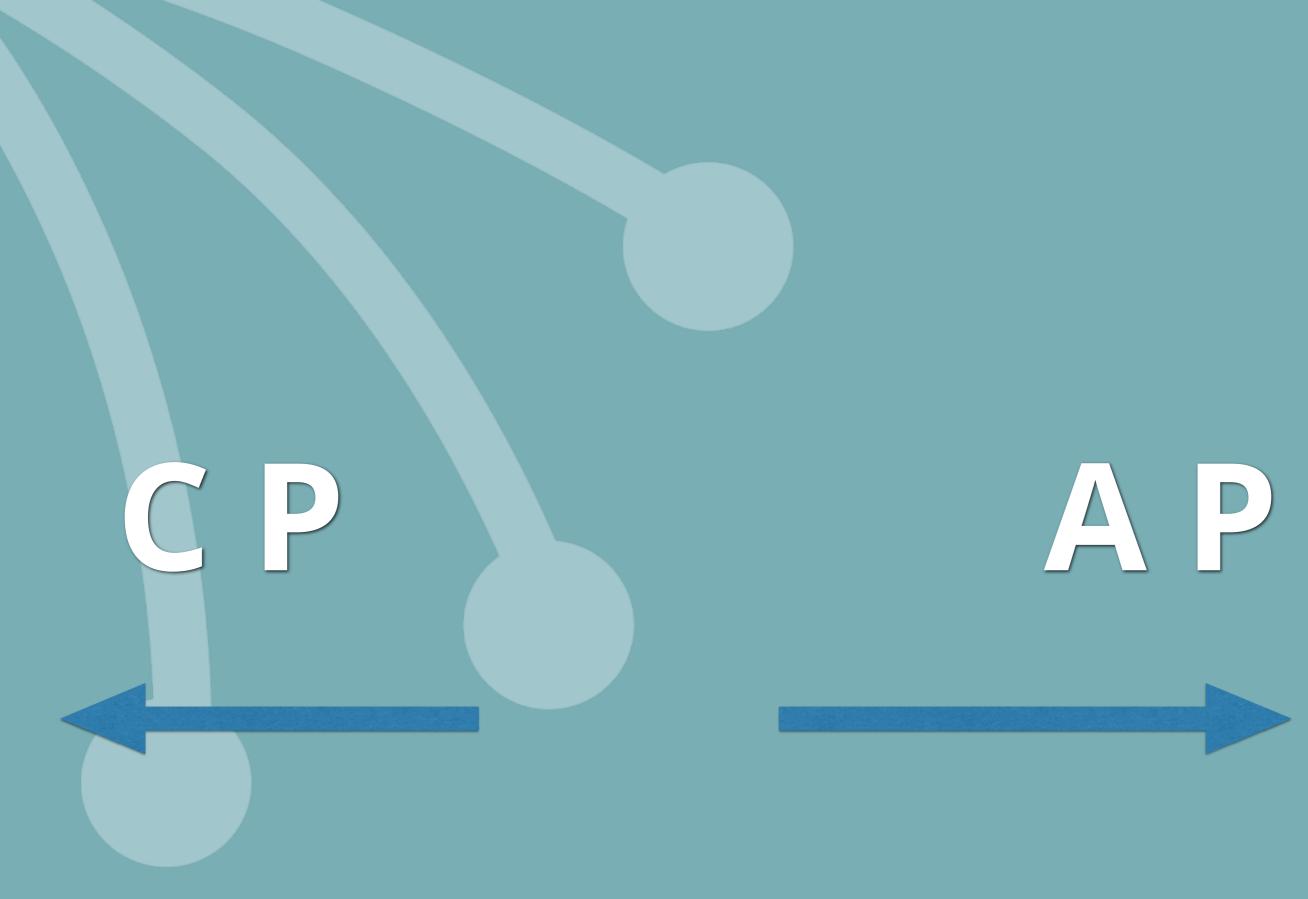


#### http://aphyr.com/posts/288-the-network-is-reliable







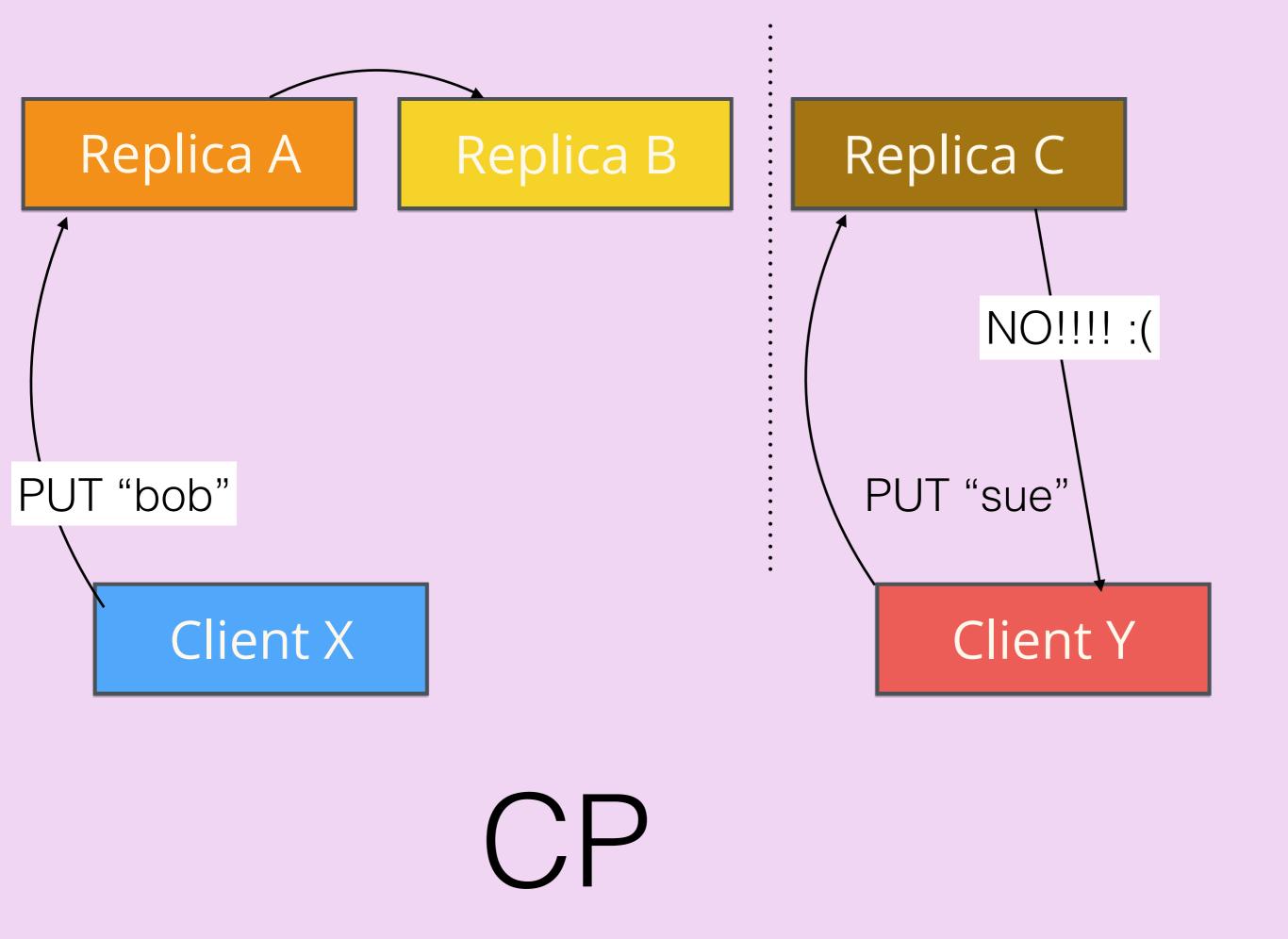


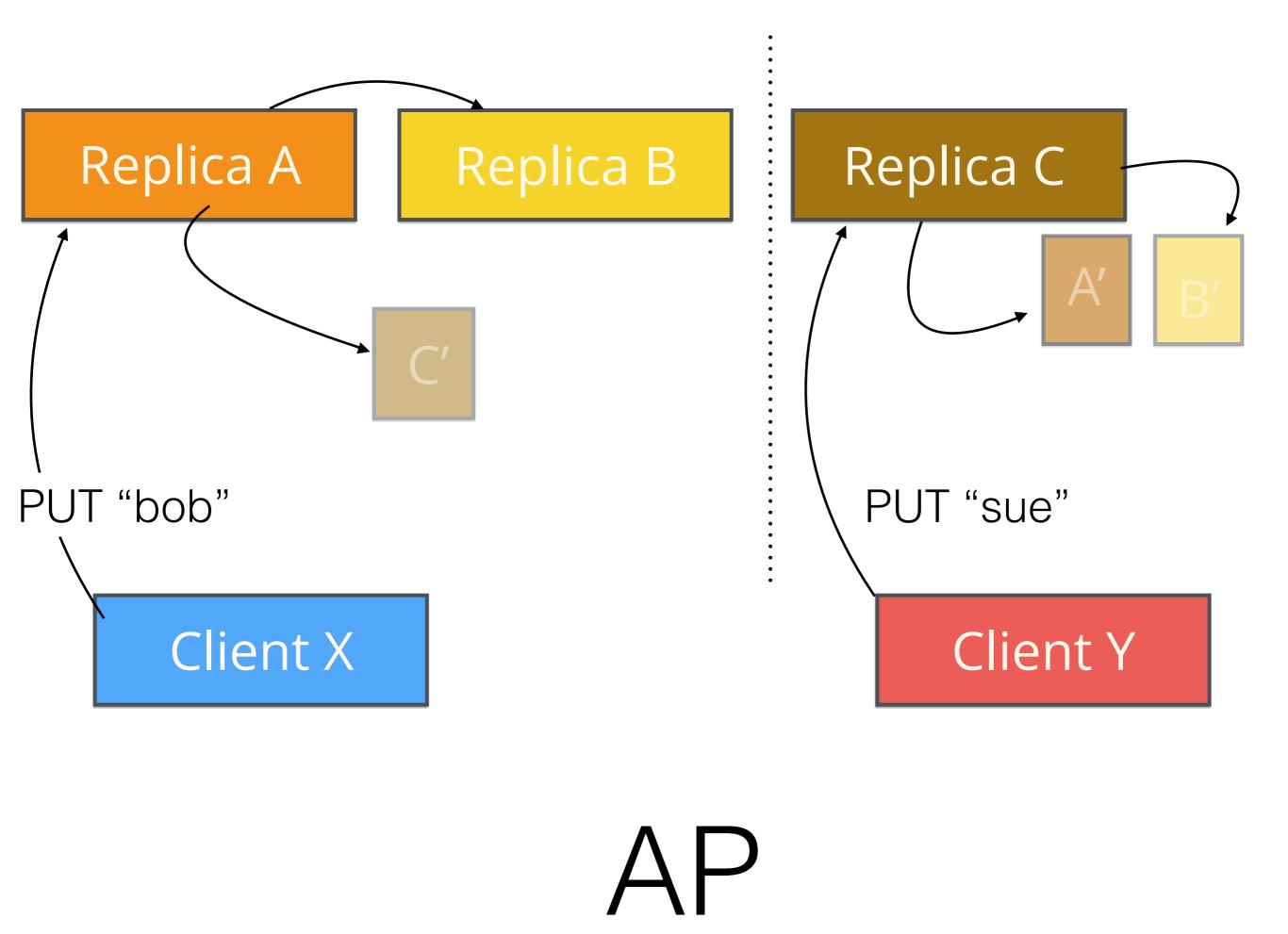
### **Eventual Consistency**

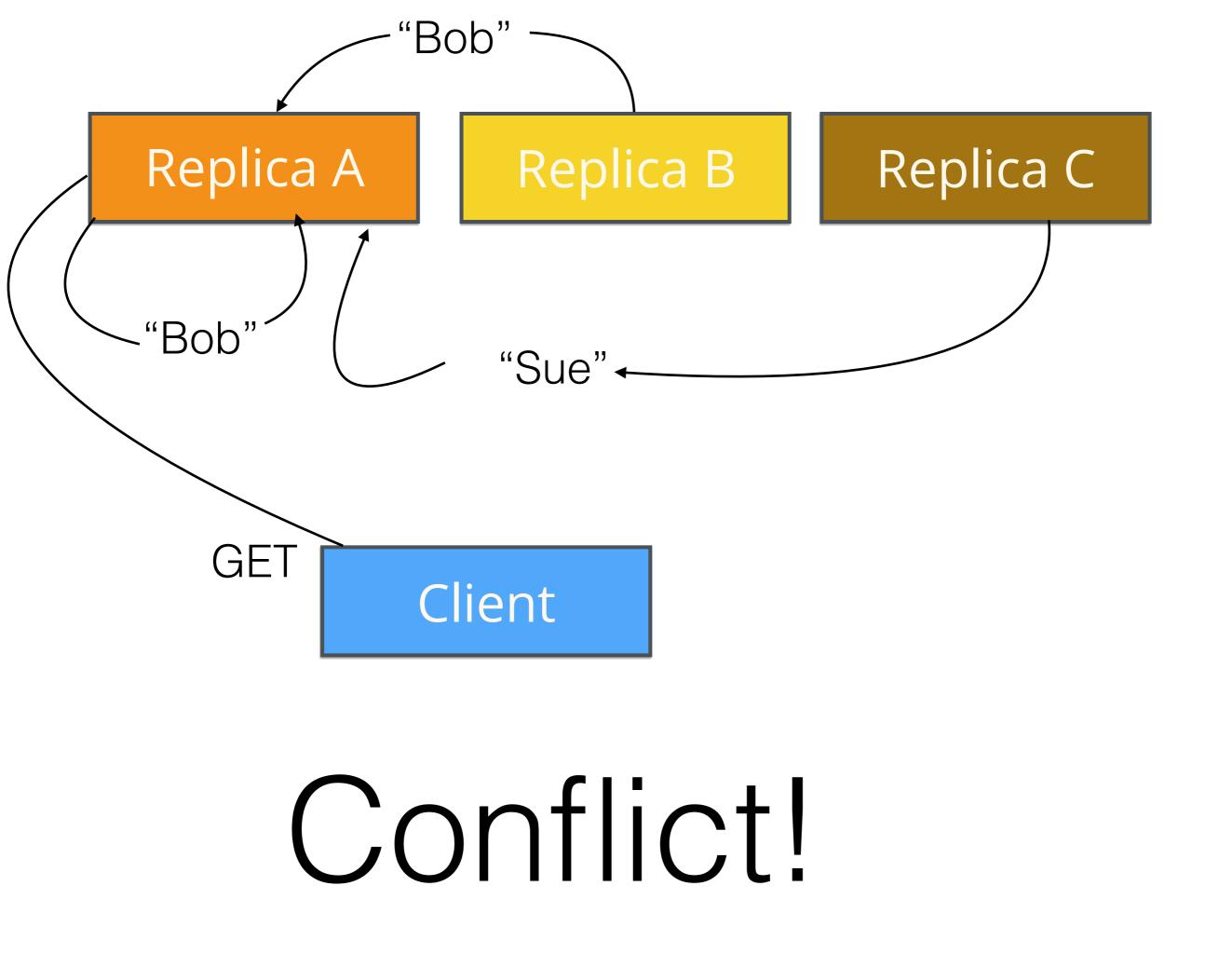
Eventual consistency is a consistency model used in distributed computing that informally guarantees that, if no new updates are made to a given data item, eventually all accesses to that item will return the last updated value.

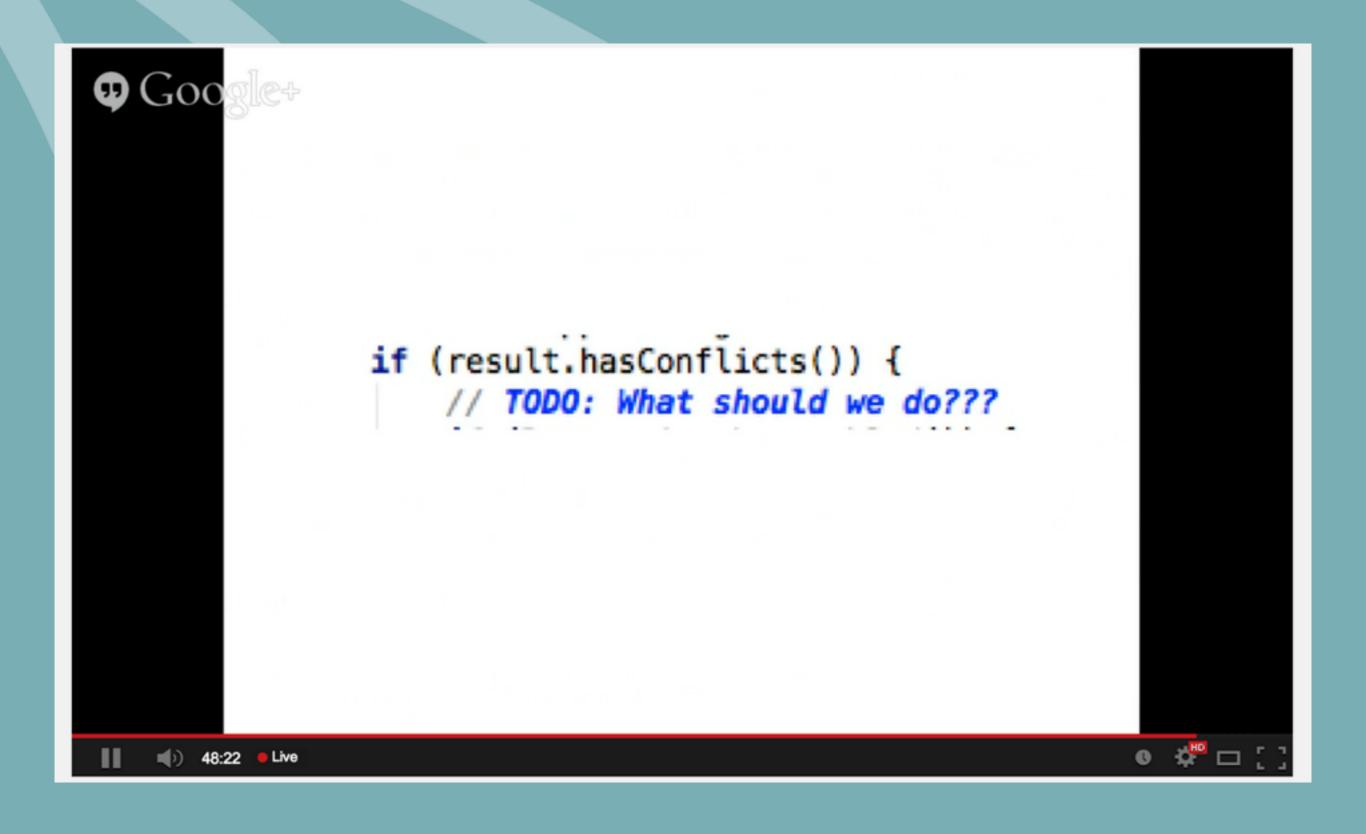
--Wikipedia

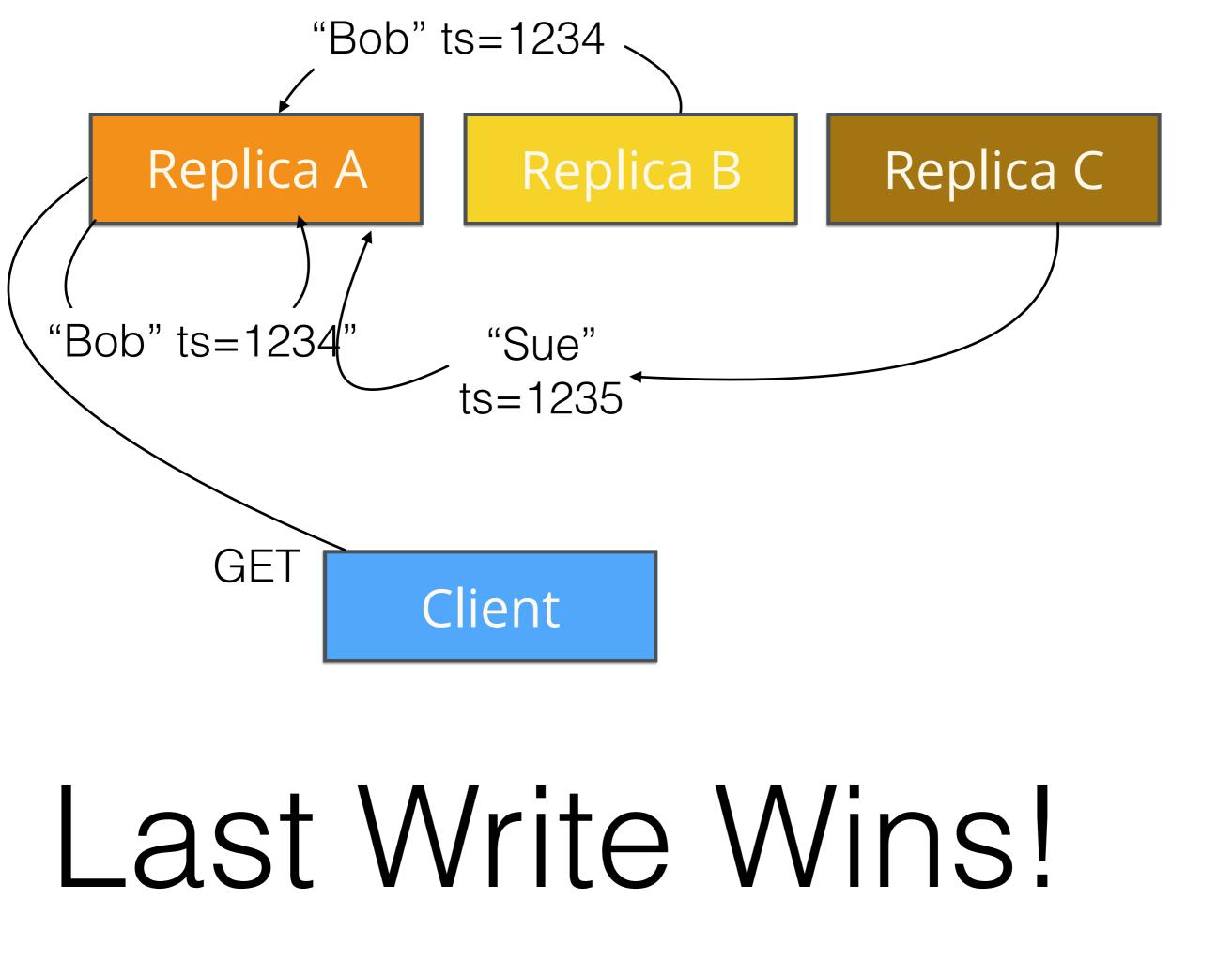


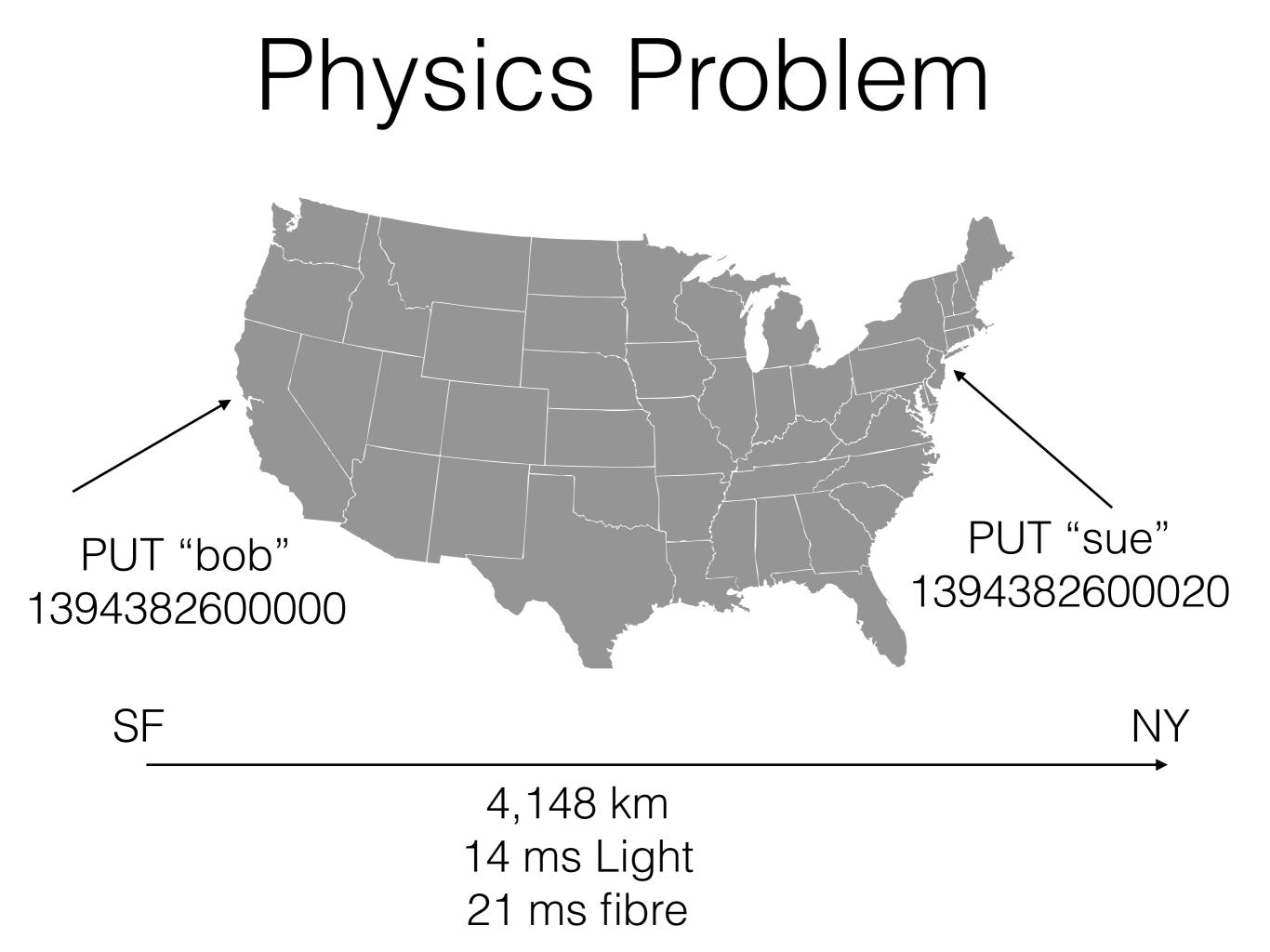


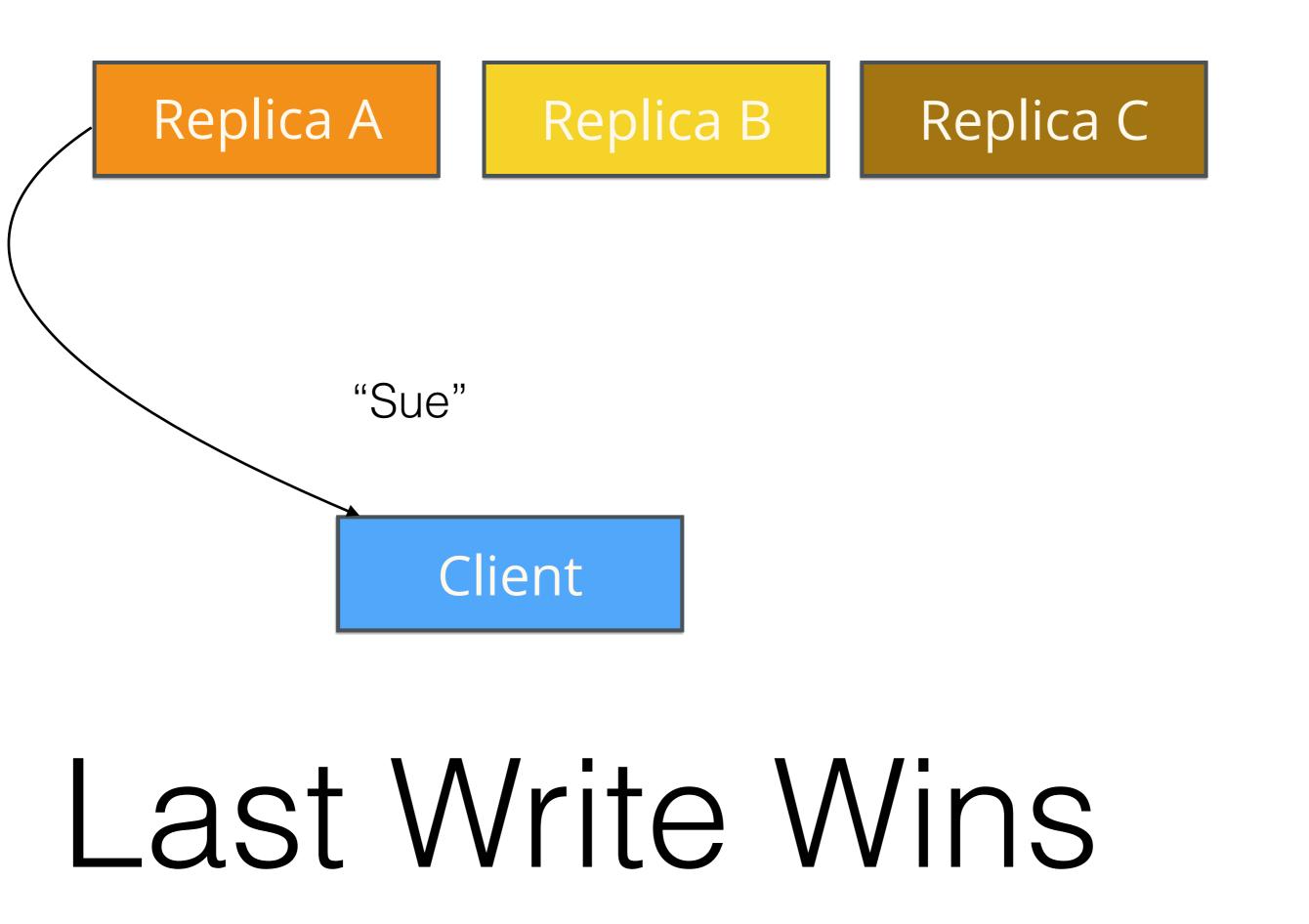


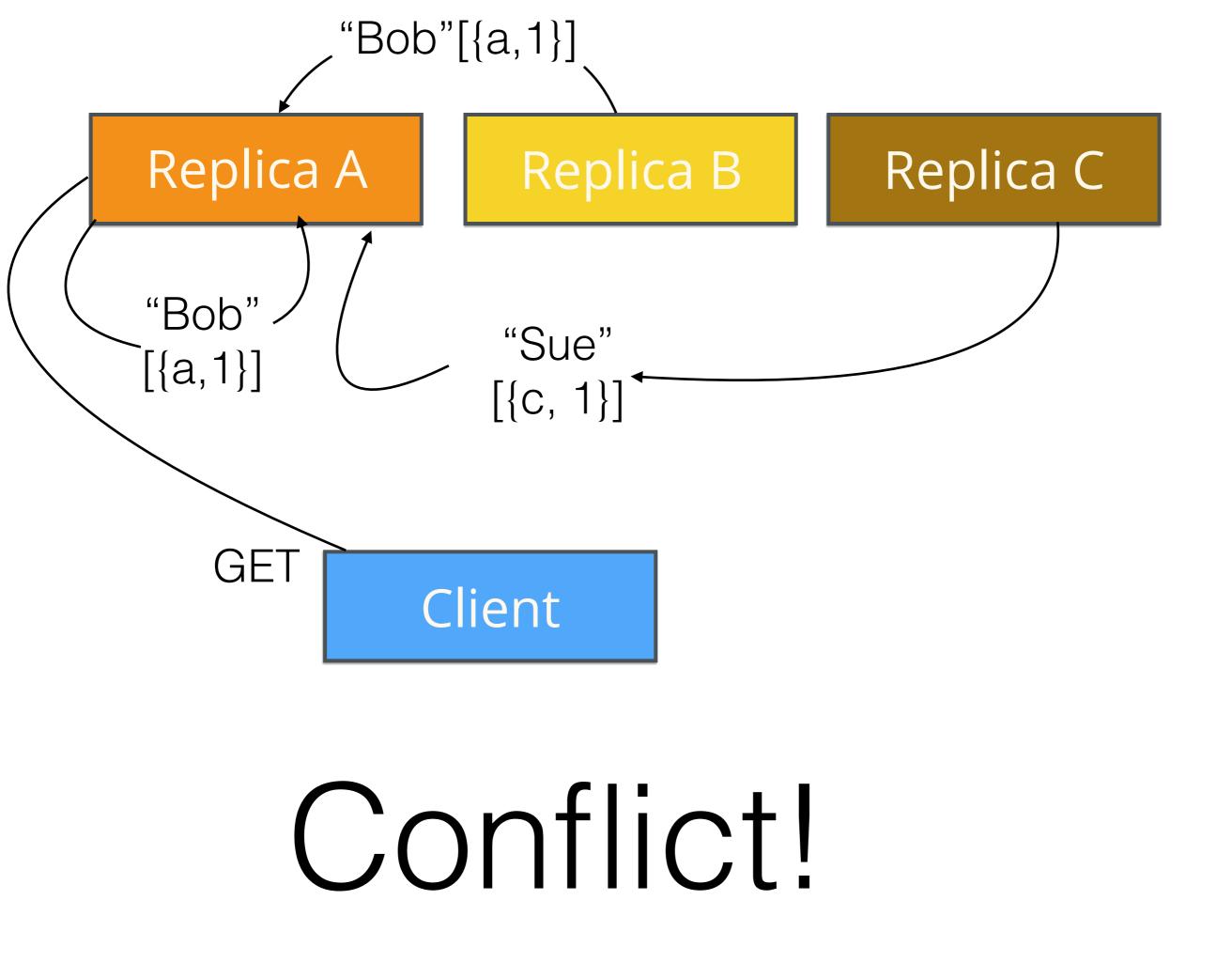


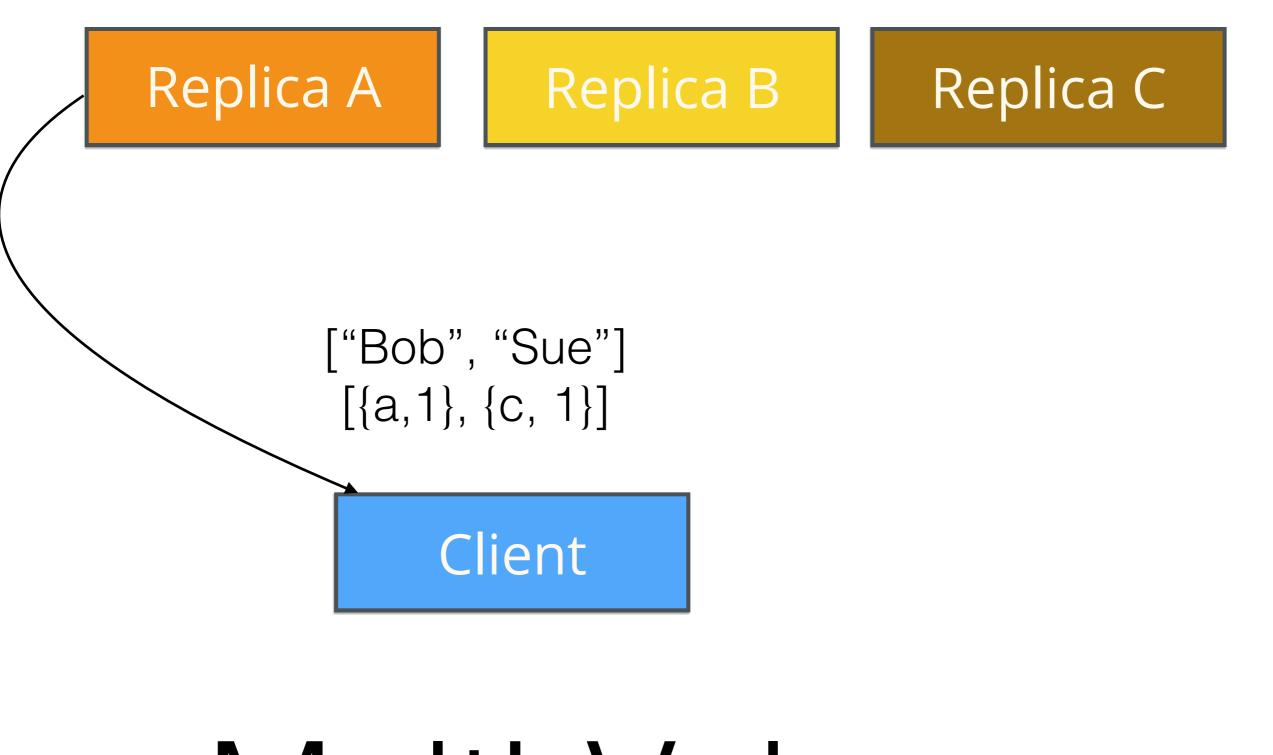








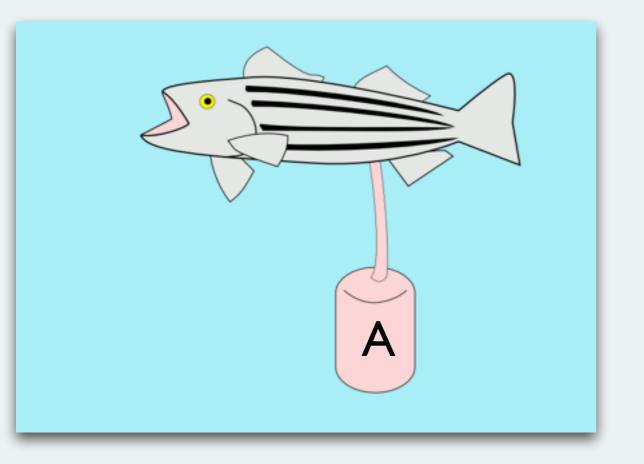


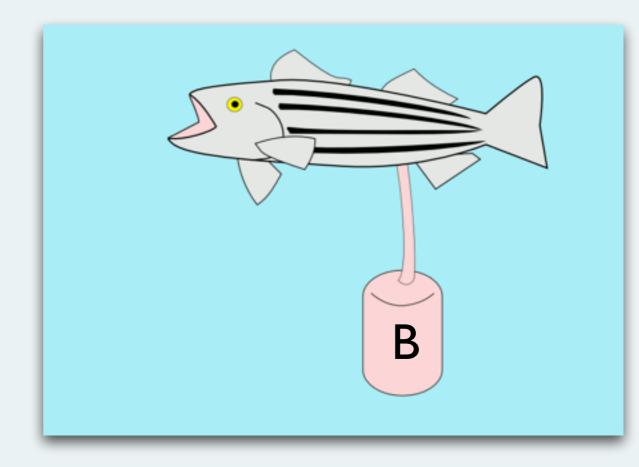


# Multi-Value

# Semantic Resolution

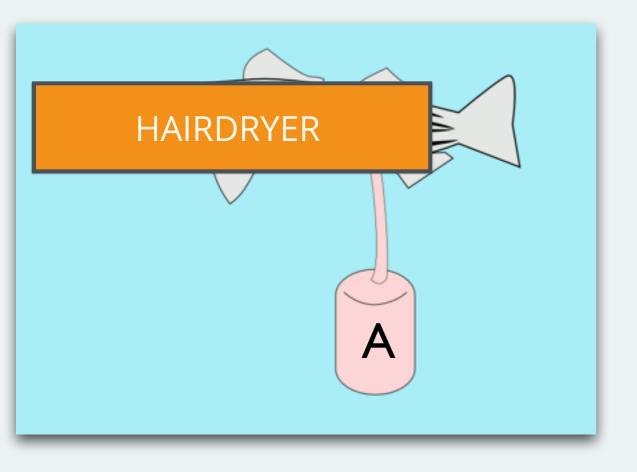
Dynamo The Shopping Cart

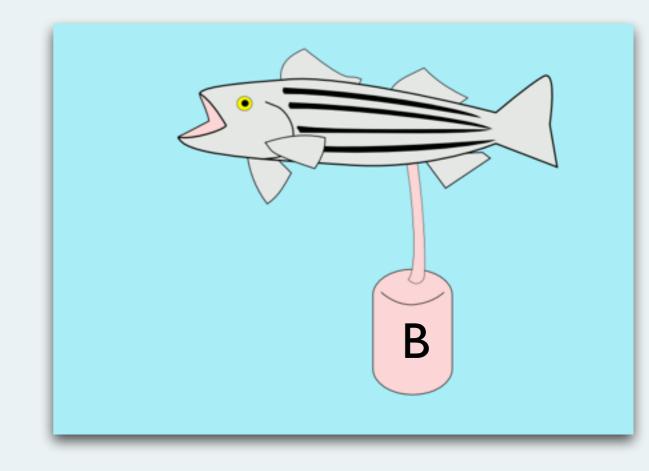




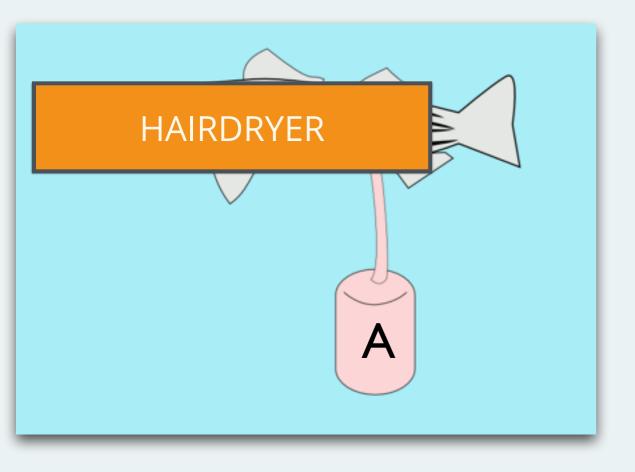
#### HAIRDRYER

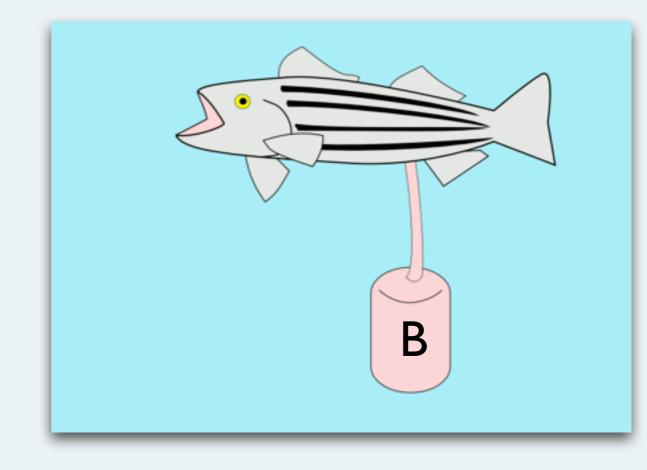






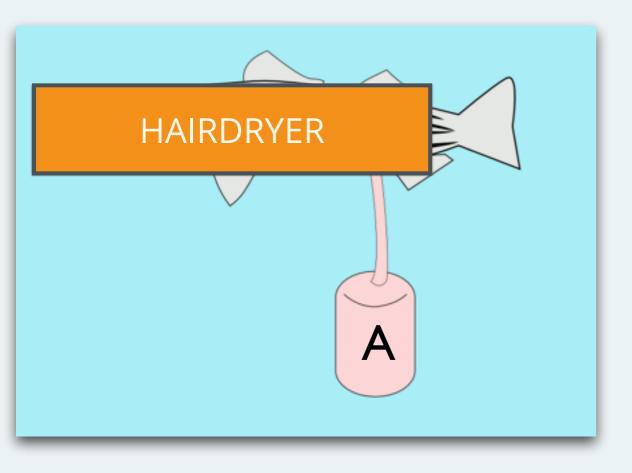


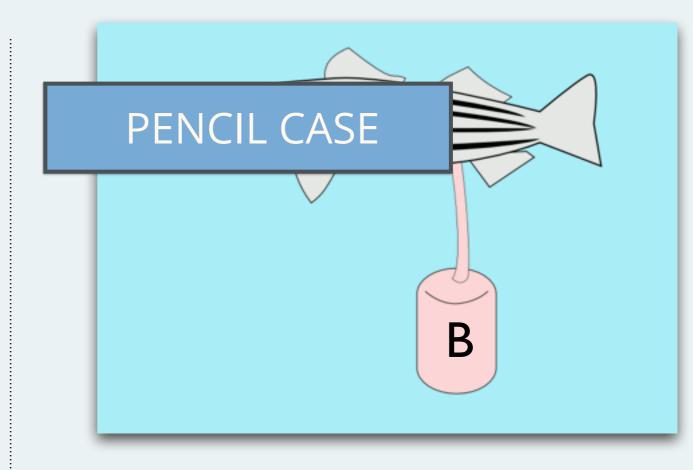




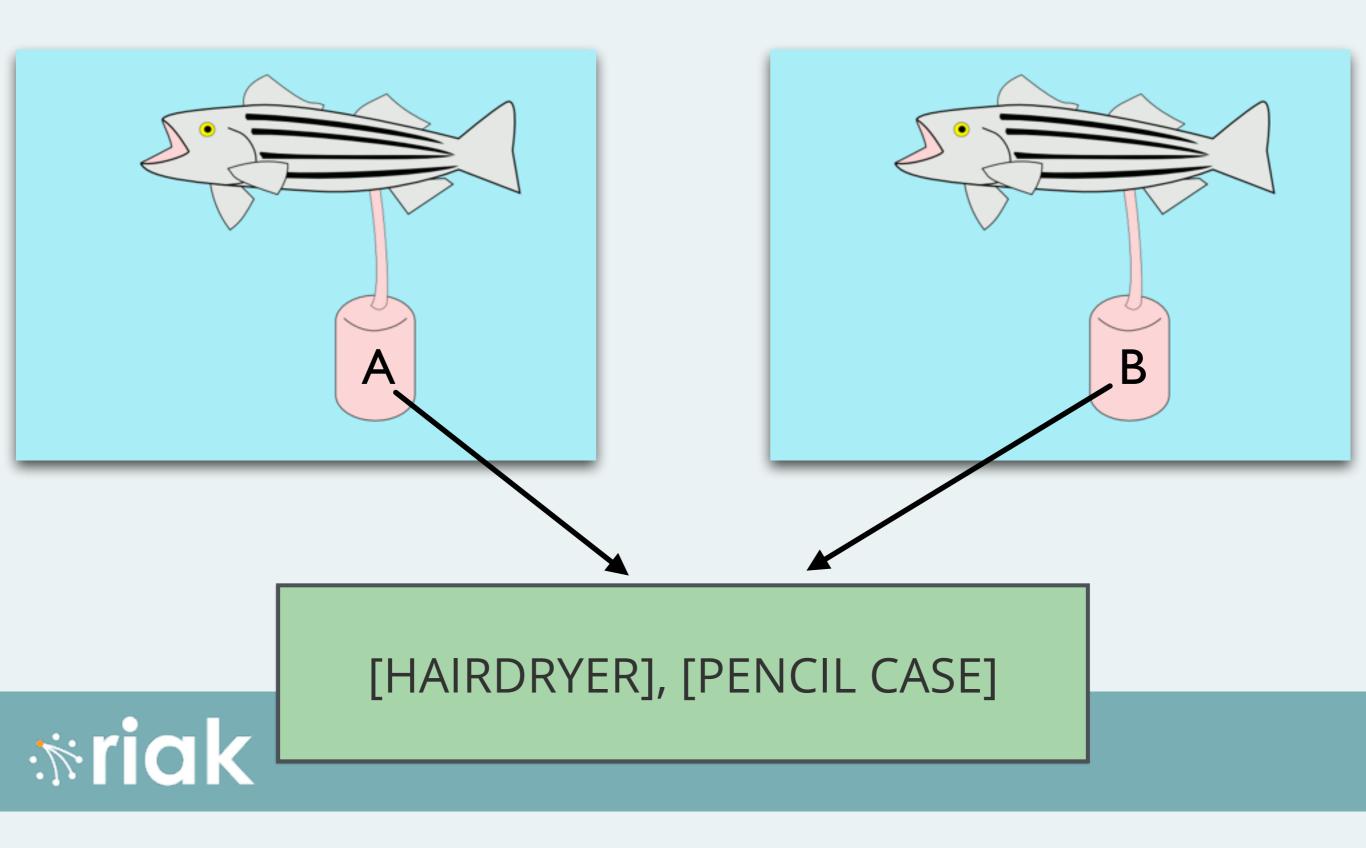












# **Nerge** Set Union of Values Simples, right?

# Removes?

Set Union? "Anomaly" Reappear

### Google F1

"We have a lot of experience with eventual consistency systems at Google."

"We find developers spend a significant fraction of their time building extremely complex and errorprone mechanisms to cope with eventual consistency"

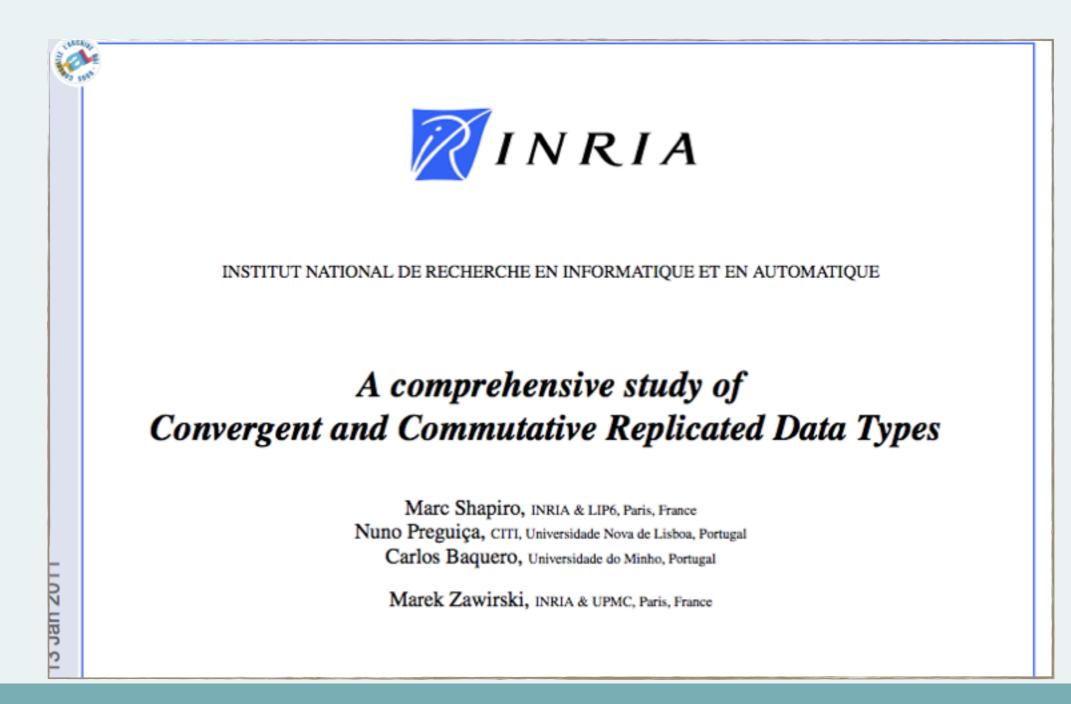


# Google F1

"Designing applications to cope with concurrency anomalies in their data is very error-prone, timeconsuming, and ultimately not worth the performance gains."











# Partially ordered set; Bottom; least upper bound $\langle S, L \rangle$



### **Join Semi-lattice** Associativity: (XUY)UZ = XU(YUZ)



### Join Semi-lattice Commutativity: XUY = YUX



#### Idempotent: X\UX = X

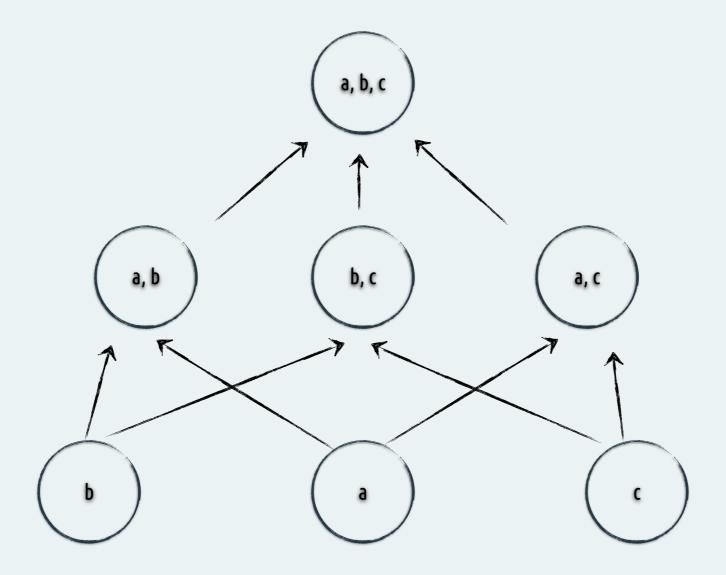


#### **Objects grow over time; merge computes LUB**

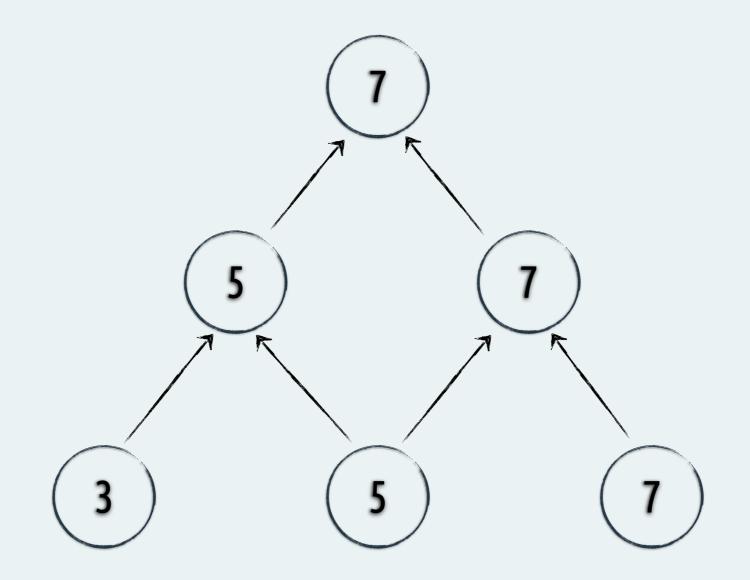


**Examples** 

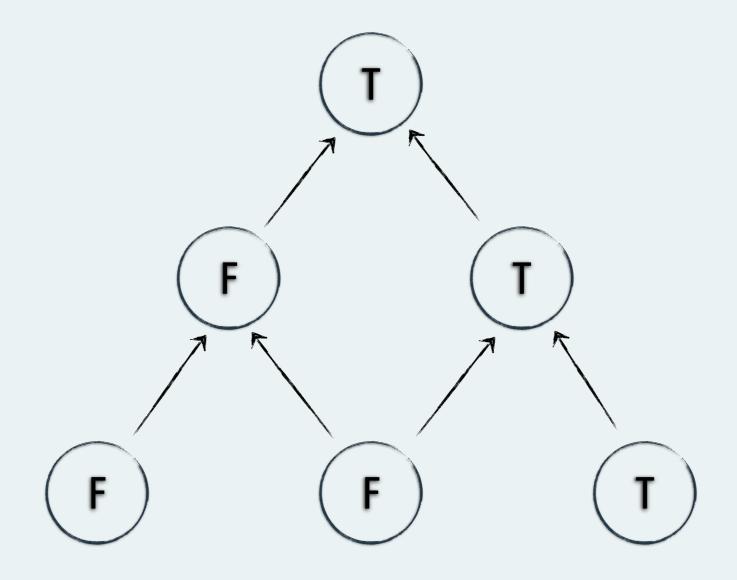




Set; merge function: union.



#### Increasing natural; merge function: max.



Booleans; merge function: or.

Deterministic Idempotent Associative Commutative

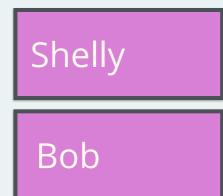
Reusable defined semantics

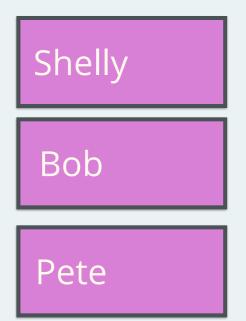
# **Evolution of a Set**

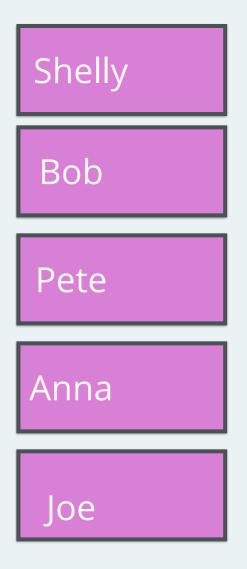


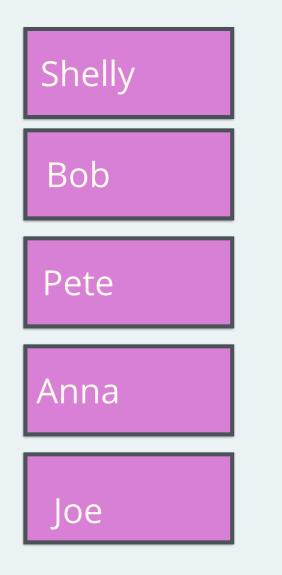
# **Evolution of a Set**



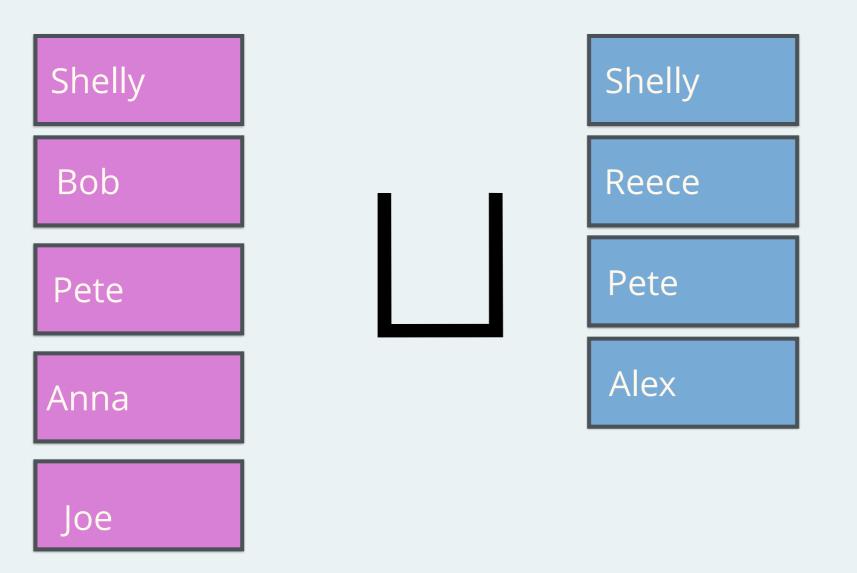


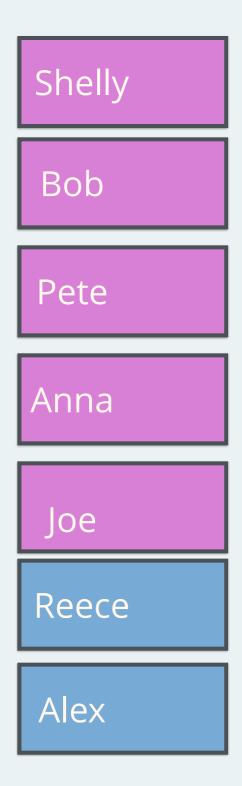




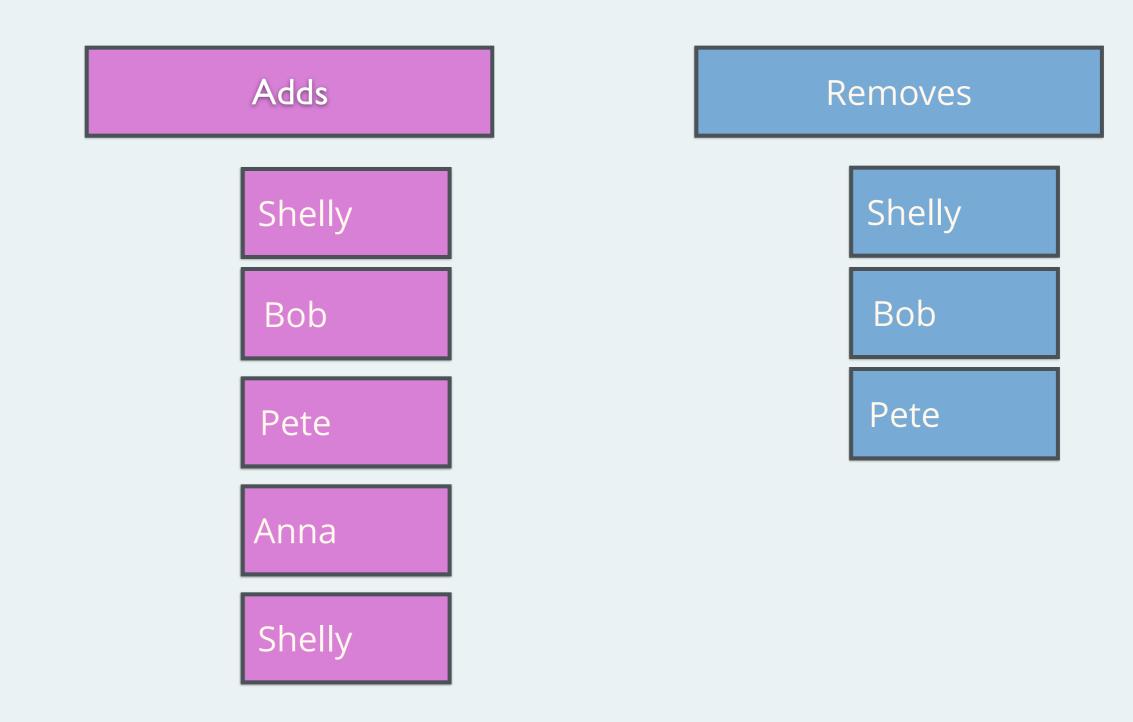


Shelly	
Reece	
Pete	
Alex	





# Evolution of a Set GISE

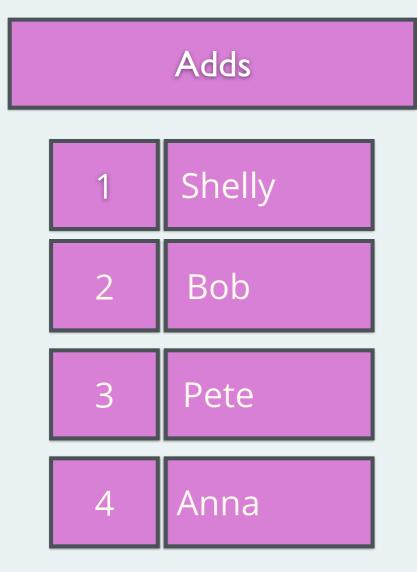


Adds	Removes	
Shelly	Shelly	
Bob	Bob	
Pete	Pete	
Anna		

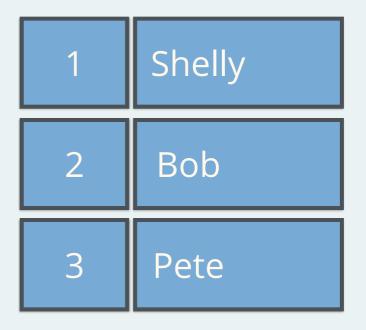
Anna

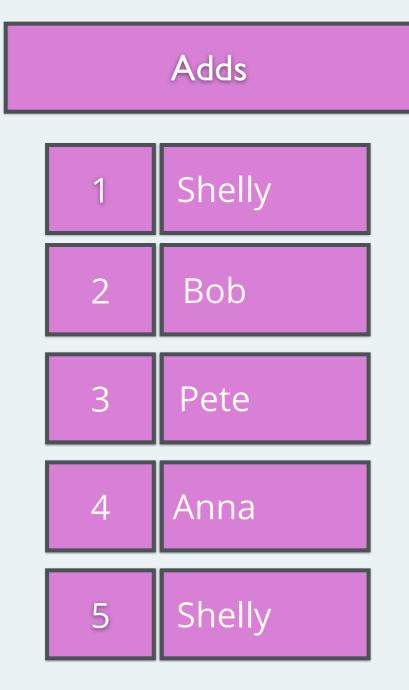
# **Evolution of a Set**

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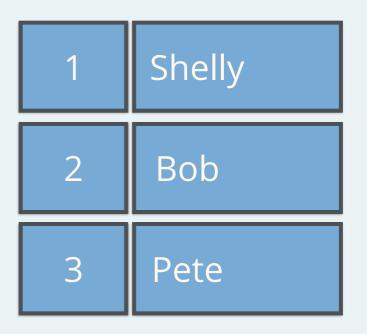


### Removes

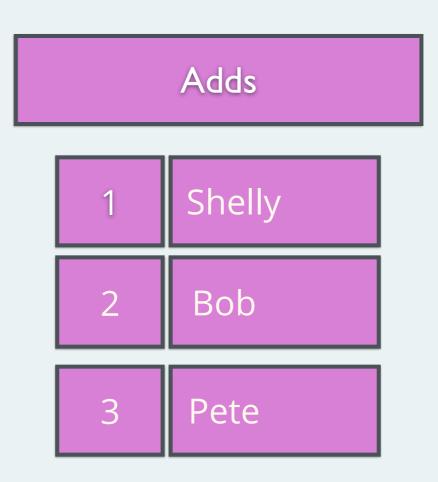




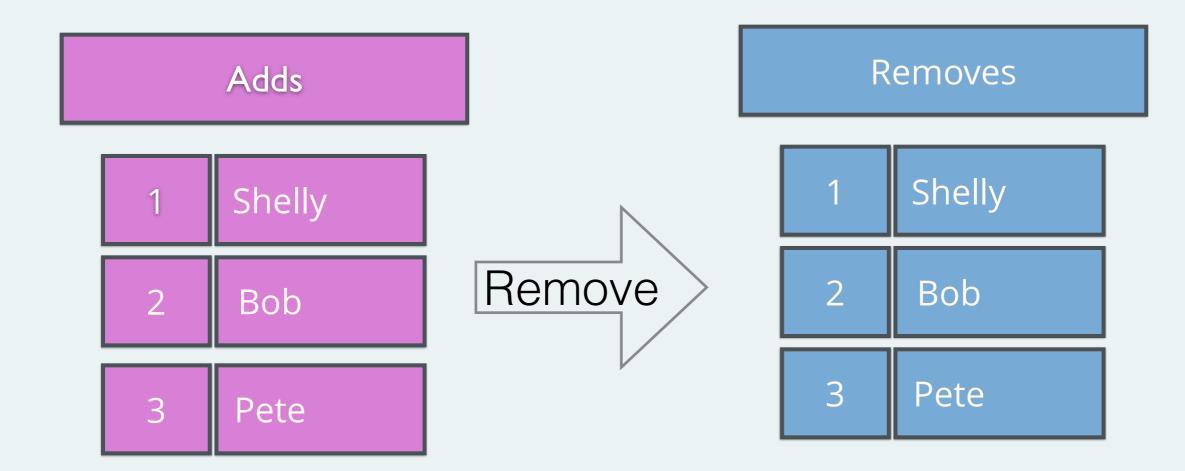
#### Removes







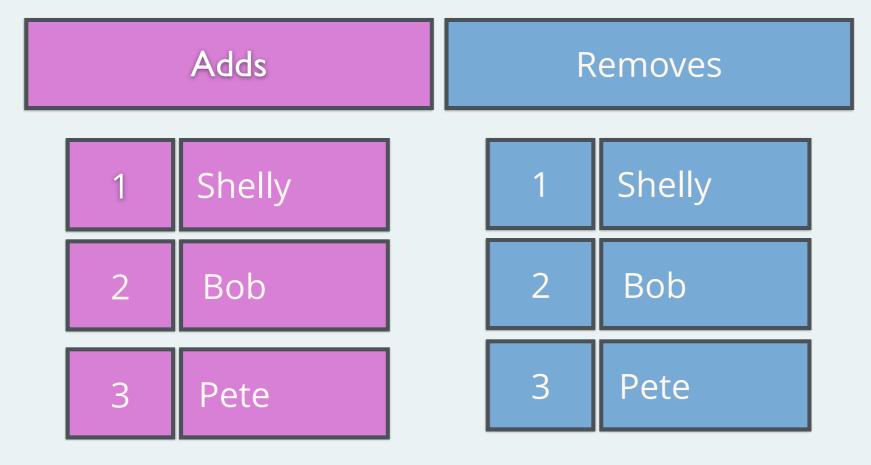




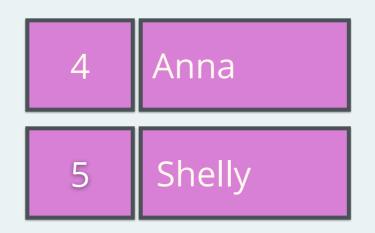


#### Adds

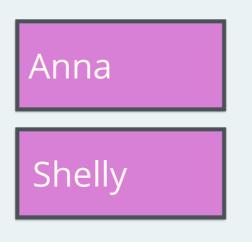
4	Anna
5	Shelly







Adds			R	emoves
1	Shelly		1	Shelly
2	Bob		2	Bob
3	Pete		3	Pete
4	Anna			
5	Shelly			



# Semantics Wins

# **Evolution of a Set**



#### **Dotted Version Vectors: Logical Clocks for Optimistic Replication**

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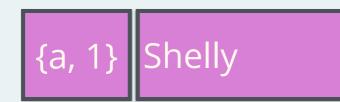
#### Abstract

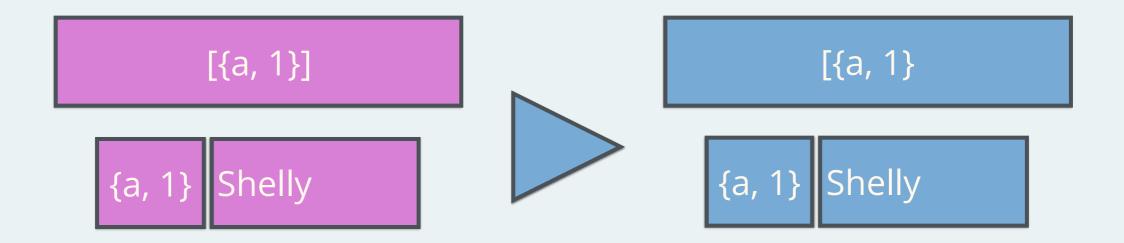
In cloud computing environments, a large number of users access data stored in highly available storage systems. To provide good performance to geographically disperse users and allow operation even in the presence of failures or network partitions, these systems often rely on optimistic replication solutions that guarantee only eventual consistency. In this scenario, it is important to be able to accurately and efficiently The mentioned systems follow a design where the data store is always writable. A consequence is that replicas of the same data item are allowed to diverge, and this divergence should later be repaired. Accurate tracking of concurrent data updates can be achieved by a careful use of well established causality tracking mechanisms [5], [6], [7], [8]. In particular, for data storage systems, version vectors [6] enables the system to compare any pair of replica versions and detect if

\*riak

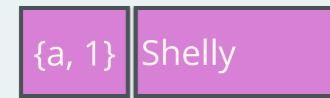
# **Evolution of a Set** U-SET OR-SET OR-SWOT

#### [{a, 1}]

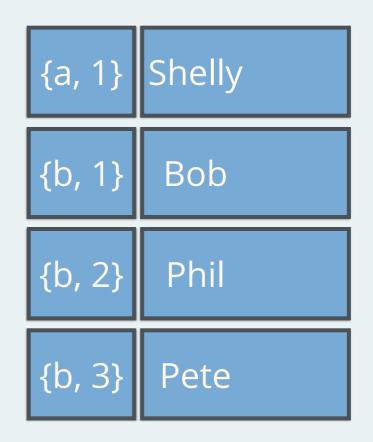


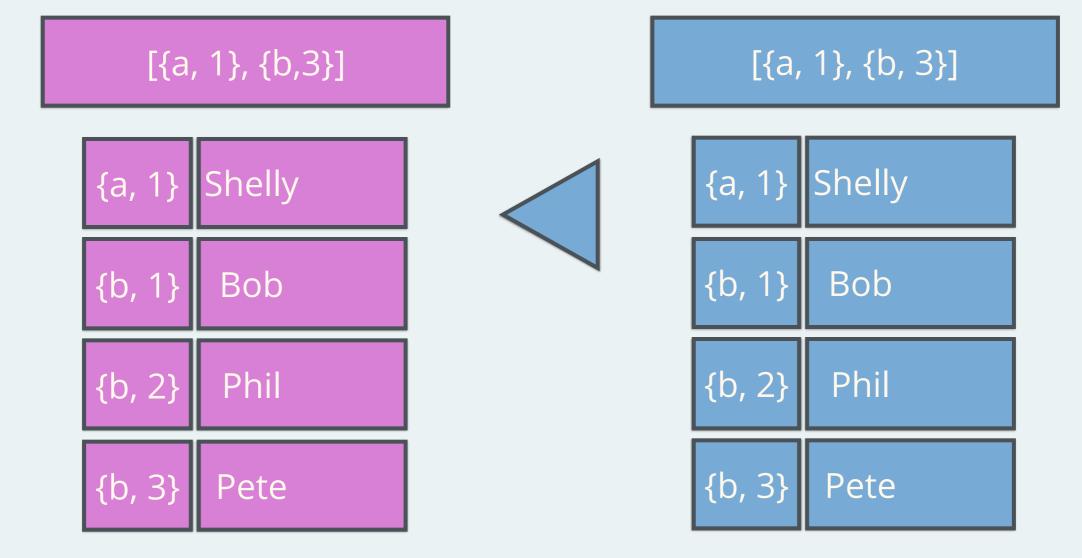


#### [{a, 1}]



#### [{a, 1}, {b, 3}]



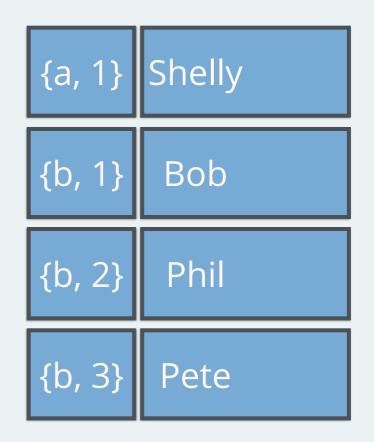


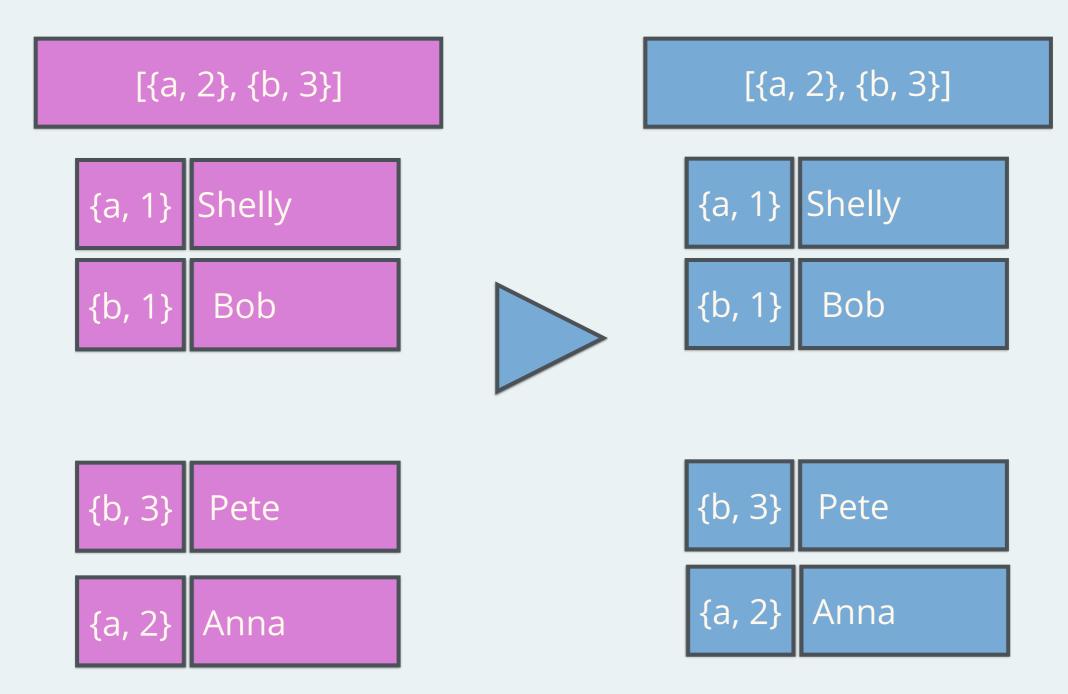
#### [{a, 2}, {b, 3}]

{a, 1}	Shelly
{b, 1}	Bob

{b, 3}	Pete
{a, 2}	Anna

#### [{a, 1}, {b, 3}]





# CRDTs

- Principled Merge
- Data Types with Defined Semantic
- Fine Grained Causality
- Building Block of EC Systems